# SGV: Spatial Graph Visualization

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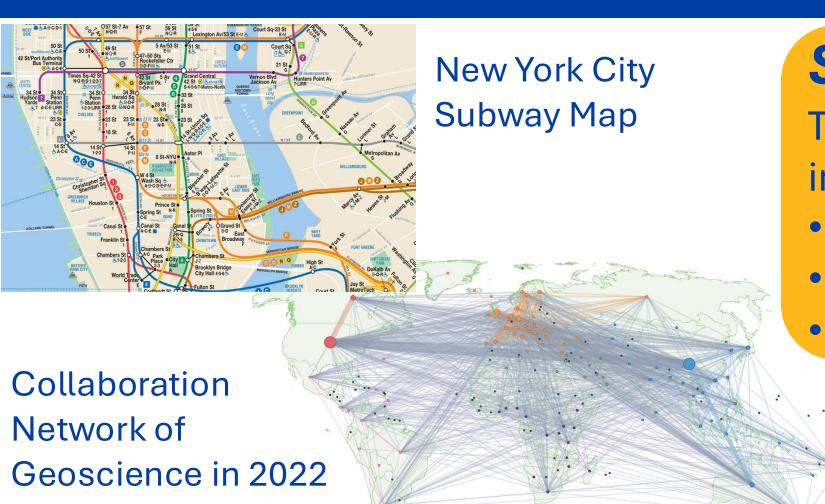
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**UCIrvine** 

# Background and Motivation



## Spatial Graph Visualization:

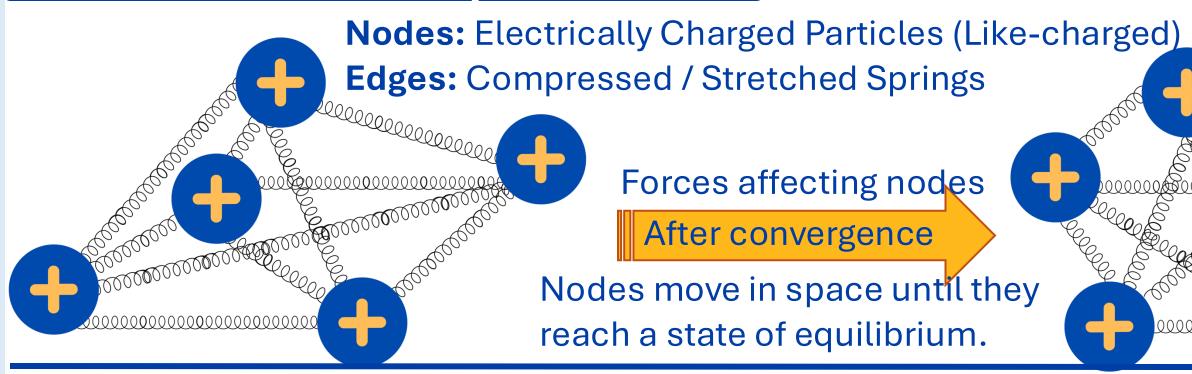
The **underlying space information** is also important. Examples:

- Railroad network
- Social Media Network
- Acaedemic Collaboration Network

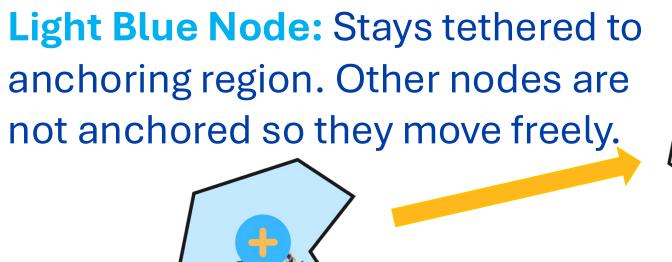
#### **Spatial Graph Model:**

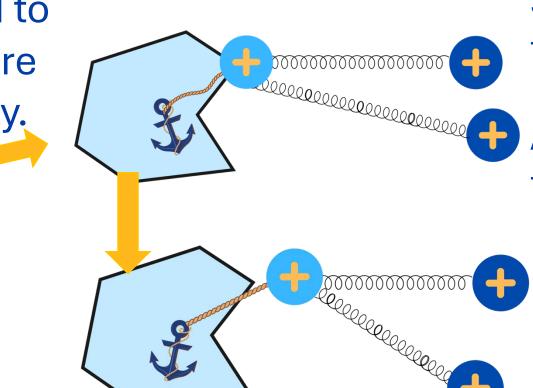
Nodes: Associated with geographic information. Edges: Undirected - Only Show Connectivity

### Force Directed Graph Model:



Forces affecting nodes After convergence Nodes move in space until they





**Spatial Graph:** Nodes have a true geographic location. In final layout, should not move too far. Anchoring force keeps nodes tethered to anchor regions.

#### **Anchor Region Types:**

- Point
- Linestring
- Polygon
- Multi-point

## Method

## Force Formulas:

#### Three types of forces between nodes:

### **Attractive (Spring) Forces:**

$$\mathbf{F}_{u \leftarrow v}^{\text{att}} = -\frac{d}{L} \, \Delta \mathbf{r}_{uv}$$

### **Repulsive Forces:**

$$d^{2} = \max(\|\mathbf{r}_{uv}\|^{2}, \varepsilon^{2}), \qquad \mathbf{F}_{u \leftarrow v}^{\text{rep}} = c_{\text{rep}} \frac{\mathbf{r}_{uv}}{d^{2}}.$$
$$\mathbf{F}_{i}^{\text{rep}} = \sum_{j \in \mathcal{N}_{r}(i)} c_{\text{rep}} \frac{\mathbf{p}_{i} - \mathbf{p}_{j}}{\max(\|\mathbf{p}_{i} - \mathbf{p}_{j}\|^{2}, \varepsilon^{2})}.$$

## **Anchoring Forces:**

$$\mu(A_u) = \begin{cases} 1, & A_u \text{ is a point,} \\ |A_u|, & A_u \text{ is a finite set of points,} \\ \text{Length}(A_u), & A_u \text{ is a linestring,} \\ \text{Area}(A_u), & A_u \text{ is a polygon,} \end{cases}$$

$$\mathbf{c}_u = \frac{1}{\mu(A_u)} \int_{A_u} \mathbf{x} \, \mathrm{d}\mu(\mathbf{x}), \quad \mathbf{p}_u = \underset{\mathbf{y} \in A_u}{\operatorname{arg\,min}} \|\mathbf{r}_u - \mathbf{y}\|.$$

### **Anchoring Force Models:**

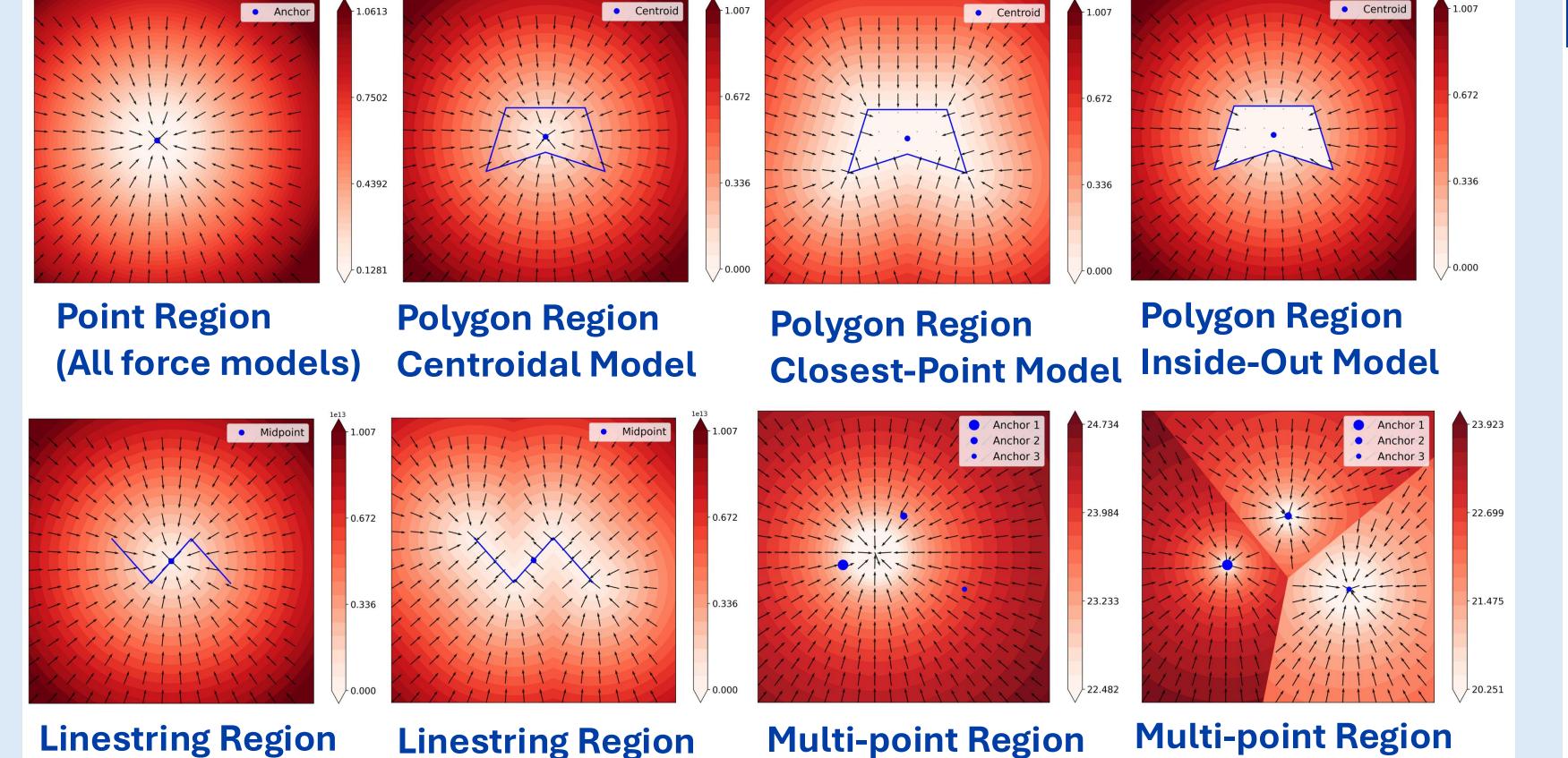
 $\mathbf{F}_{u}^{\text{anchor}} = \alpha \left( \mathbf{c}_{u} - \mathbf{r}_{u} \right).$ Centroidal:

Inside-Out:

**Closest-Point:** 

**Closest-Point Model** 

# **Anchoring Force Fields:**



## Scalability:

**Centroidal Model** 

### Problem: Spatial Graphs can grow arbitrarily large

Naïve Calculation of Repulsive forces between all nodes:  $O(N^2) \rightarrow Bottleneck$ **Spark SQL:** Parallelize the computations in distributed form  $\rightarrow$  Increases Scalability **Future Work:** 

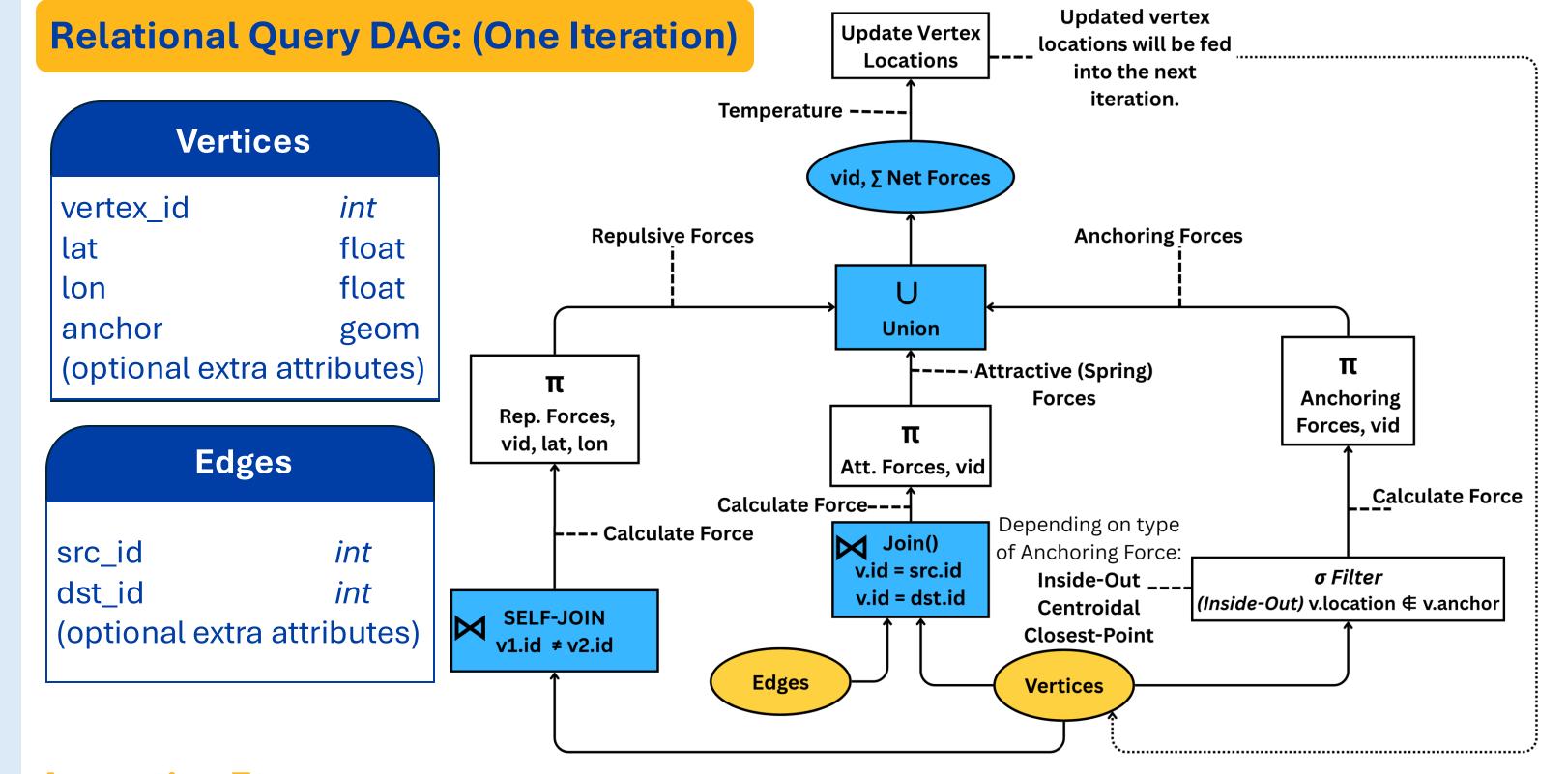
**Closest-Point Model Centroidal Model** 

Use a Spatial Index (QuadTree partitioning of the space) and approximate effect of nodes that are further than the distance threshold theta. Based on well-separated pairs.

## **Related Work:**

Polygonally Anchored Graph Drawing (Extended Abstract) (Chiu et al – 2024) Well Separated Pair Decomposition (Chan - 2008) A Potential-Field-Based Multilevel Algorithm for Drawing Large Graphs. (Hachul – 2005)

# Relational Query Modeling



#### **Attractive Forces:**

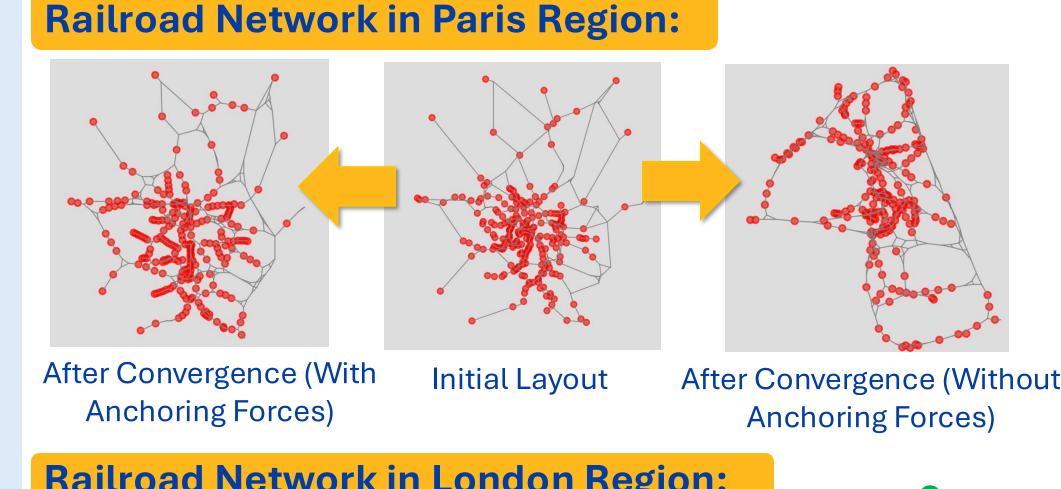
JOIN: Edges ⋈ Vertices → Spring force alpha L (ideal spring length) applied to both dst and src vertices. **Repulsive Forces:** 

**SPATIAL SELF-JOIN: Vertices** ⋈ **Vertices** (within theta) → Repulsive force applied to both vertices. **Anchoring Forces:** 

**VERTICES** → pull-to-anchor (centroid / closest-point) (σ Filter inside-out: skip if inside region) → per-vertex f **Net Force Aggregation:** 

U UNION ALL (attractive, repulsive, anchoring)  $\rightarrow \Sigma$  by vid  $\rightarrow F_{total} \rightarrow Updater$ 

## Use Cases



Railroad Network in London Region: A) Initial Layout B) After Convergence (With **Anchoring Forces**)

Green **Nodes:** Main stations **Strong Anchoring** Forces → Zoomed in – Railroad network in Nodes stay tethered to their true

Paris: Before (Top) vs After (Bottom) Convergence – Using Anchoring Forces. More homogenous edge lengths and increased layout legibility

## Results and Conclusion

geographic

locations.

**Evaluation Metrics:** 

Higher value is desired

**HEL: (Homogenous Edge Lengths)** 

## **Experiment Setup:**

We conducted experiments on a 12-node Spark cluster. The master node has two 8-core Intel Xeon E5-2609 v4 (1.7 GHz) CPUs and 128 GB RAM; each worker has two 6-core Xeon E5-2603 v4 CPUs and 64 GB RAM. All nodes run CentOS 7.5 with local SSDs for the OS and HDDs for data.

Dataset	# Vertices	# Edges
Gowalla Subset	22,803	381,384
Gowalla	107,092	913,660
Gowalla x2	214,184	1,827,320
Gowalla x5	535,460	4,568,300
Gowalla x10	1,070,920	9,136,600
Gowalla x 100	10,709,200	91,366,000

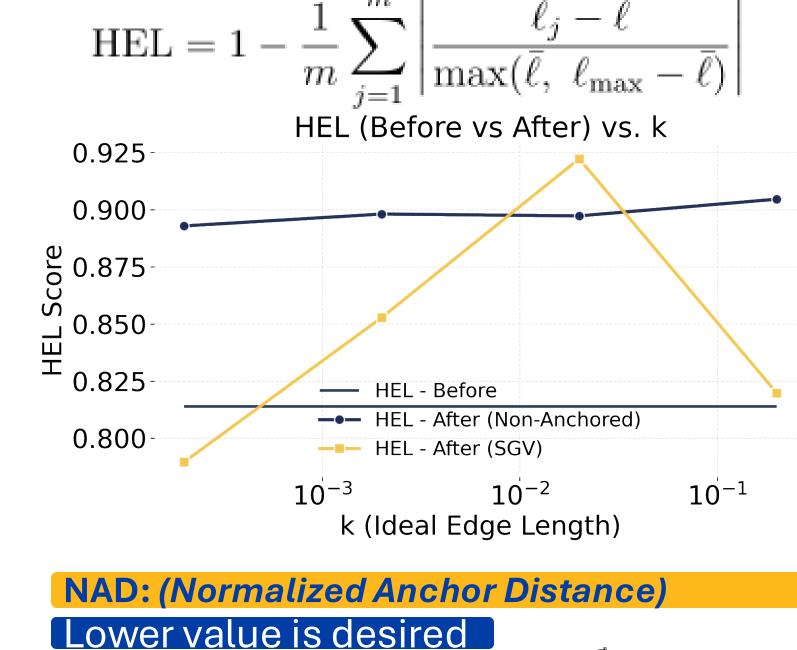
Goal: find a balance between edge length homogeneity and average distance from anchor of nodes in final layout.

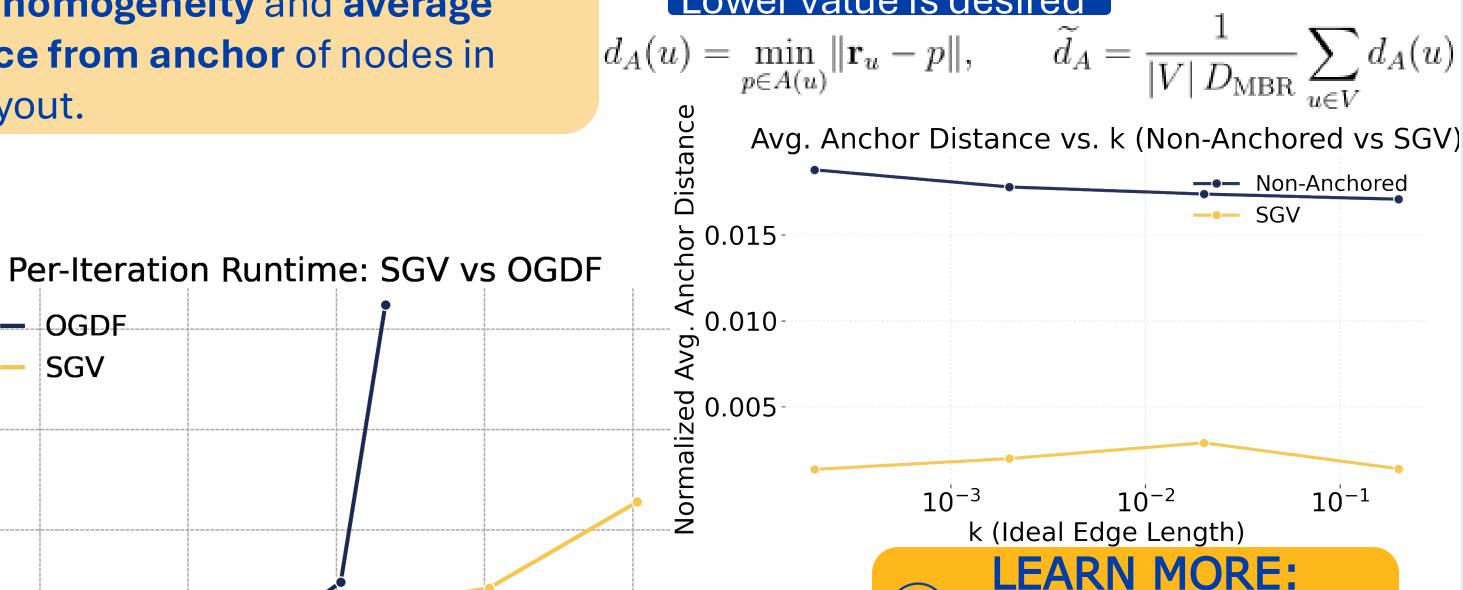
--- OGDF

 $10^{3}$ 

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SGV





github.com/tarlaun/fdgv

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# of Vertices This work is supported in part by the National Science Foundation (NSF) under grant IIS-2046236

 $10^6$